

Question 1

Recall from lecture that we defined a *model* of DPDL to be a set X equipped with partial functions $\llbracket \pi \rrbracket : X \rightarrow X$ for each $\pi \in \Pi$ (where Π was our set of programs), and for each propositional letter p , a set $\llbracket p \rrbracket \subseteq X$, the *extension* of p , or the set of “all those states $x \in X$ where p is true”. We then recursively defined the extensions for various logical connectives:

$$\begin{aligned} x \in \llbracket \neg\varphi \rrbracket &\iff x \notin \llbracket \varphi \rrbracket \\ x \in \llbracket \varphi \wedge \psi \rrbracket &\iff x \in \llbracket \varphi \rrbracket \cap \llbracket \psi \rrbracket \\ x \in \llbracket \varphi \rightarrow \psi \rrbracket &\iff x \notin \llbracket \varphi \rrbracket \text{ or } x \in \llbracket \psi \rrbracket \\ x \in \llbracket [\pi]\varphi \rrbracket &\iff \llbracket \pi \rrbracket(x) \in \llbracket \varphi \rrbracket \text{ or } \llbracket \pi \rrbracket(x) \text{ is undefined} \\ x \in \llbracket \langle \pi \rangle \varphi \rrbracket &\iff x \in \llbracket \neg[\pi]\neg\varphi \rrbracket \end{aligned}$$

So: what has to be true about the partial function $\llbracket \pi \rrbracket$ at a state x in order for it to be the case that $x \in \llbracket [\pi]\varphi \rrbracket$ but $x \notin \llbracket \langle \pi \rangle \varphi \rrbracket$? Why?

Question 2

Annotate the following piece of C-like code with the appropriate Hoare Logic annotations to prove that it correctly computes $n!$. Assume n is given some integer value (but put in the correct precondition!). Like the example done in lecture, the final $\{ \}$ doesn't need to literally say that $\text{res} = n!$, but it should clearly imply it.

```
{
}
i:=n;
res:=1;
{
}
while (i>0) do (
  {
}
  res := res * i;
  i := i-1
  {
}
)
{
}
```
